

Operation Imaginary II

Editorial

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TW222I

Problem Statement

You have a ruler and a compass.

Draw a N -units segment, given two points separated by 1 unit.

Trivial Solution

Induction!

$O(N)$

Draw a line joining initial

Repeatedly draw circle ...

$C = 2000$

Score: 20.000

Slightly optimized ($C \sim 130$)

Set B = block size. We find $0 \dots 2B - 1$, then extend it by B each time.

$B = 25: C = 129$

$B = 30: C = 125$

$B = 31: C = 125$

$B = 32: C = 125$

$B = 33: C = 125$

$B = 34: C = 125$

$B = 35: C = 126$

$B = 40: C = 129$

$B = 45: C = 133$

$B = 50: C = 139$

$2B + N / B \sim 2B^2 = N \Rightarrow B \sim 32$

Score: 48.000 ($C = 130$)

Slightly optimized ($C = 125$)

$C = 125$ for this approach

Score: 49.298

Further optimized (C = 123)

We can skip the later part of first step if it is not used anyways

C = 32: 123

Score: 49.817

Binary! (C = 22)

Why not?

Repeat the block...

Define $\text{plot}(a, b)$: plot $2b - a$ from a and b

0, 1 $\rightarrow$ 2

1, 2 $\rightarrow$ 3

1, 3 $\rightarrow$ 5

3, 5 $\rightarrow$ 7

C = 22

Score: 81

10111

01111

00111

00101

00011

00010

00001

00000

$\sim O(2^{\text{\#bits} = 1}) = 2 * 11 = 22 \text{ (N = 1023)}$

Further into the hole

Is there any solution better than $C = 22$?

‘Any points’

Is there any exploit on this?

Hard bound ($C \geq 12$)

Note that theoretical minimum = $1 + \text{ceil}(\lg N) = 12$ ($N = 2000$)

So the possibility is not large

Further into the hole ($C = 2I$)

10111

01111

00111

00101

00011

00010

00001

00000

Note that this is maximum when binary of N is mostly consists of 1

Further into the hole ($C = 21$)

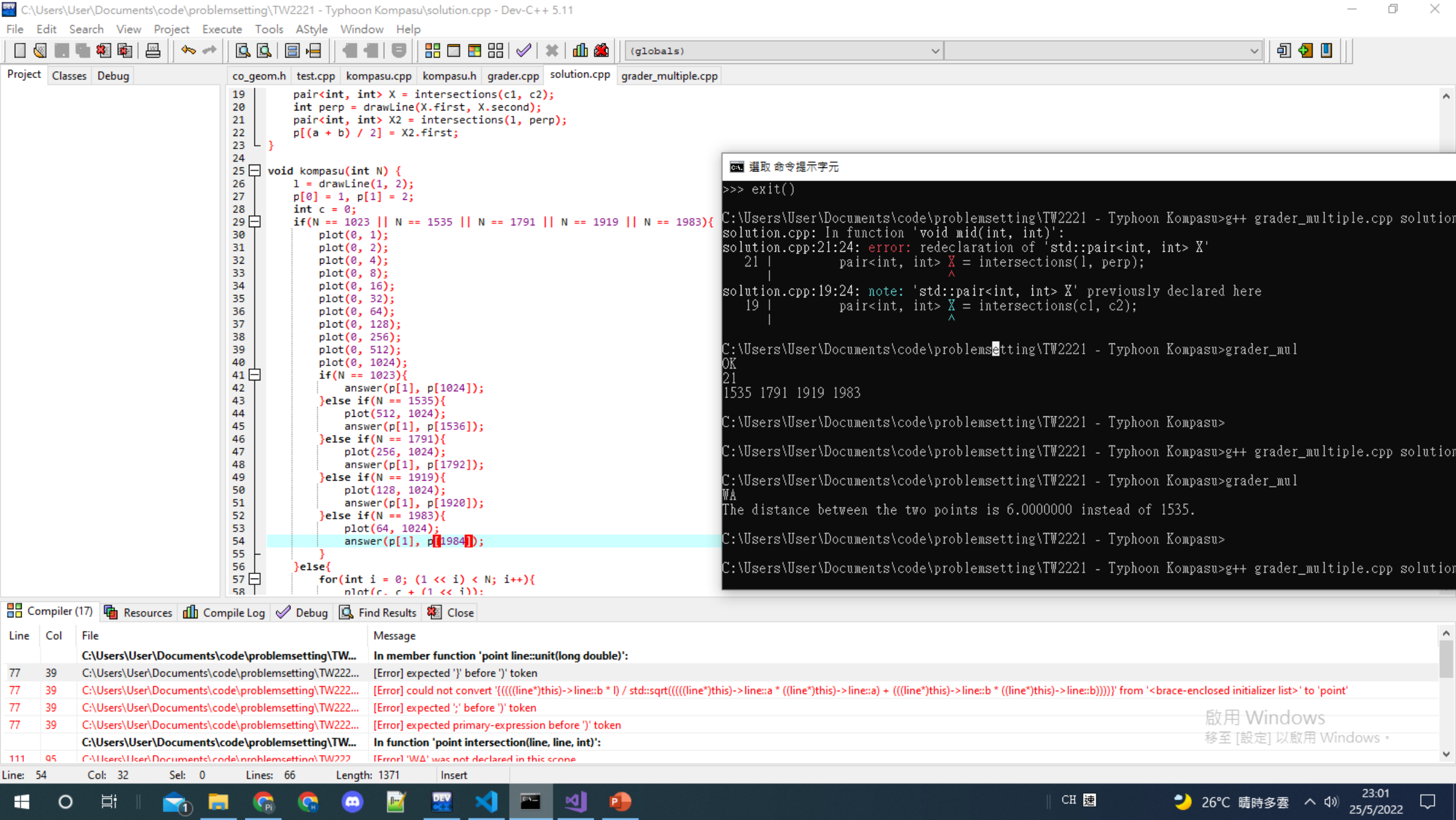
Hardcode for $N = 1023$:

0, 1, 2, 4, ..., 1024

Answer(1, 1024)

$C = 21$

Score: 82.5



Further (C = 20)

Hardcode for these cases with many 'I's!

N = 1535, 1791, 1919, 1983 (bounded by N = 2000, or actually more)

Note that just use (I, 2^M , 2^N) is sufficient

C = 20

Score: 84

C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu\solution.cpp - Dev-C++ 5.11

File Edit Search View Project Execute Tools AStyle Window Help

Project Classes Debug

co_geom.h test.cpp kompasu.cpp kompasu.h grader.cpp solution.cpp grader_multiple.cpp

```
19 pair<int, int> X = intersections(c1, c2);
20 int perp = drawLine(X.first, X.second);
21 pair<int, int> X2 = intersections(l, perp);
22 p[(a + b) / 2] = X2.first;
23 }
24
25 void kompasu(int N) {
26     l = drawLine(1, 2);
27     p[0] = 1, p[1] = 2;
28     int c = 0;
29     if(N == 1023 || N == 1535 || N == 1791 || N == 1919 || N == 1983){
30         plot(0, 1);
31         plot(0, 2);
32         plot(0, 4);
33         plot(0, 8);
34         plot(0, 16);
35         plot(0, 32);
36         plot(0, 64);
37         plot(0, 128);
38         plot(0, 256);
39         plot(0, 512);
40         plot(0, 1024);
41         if(N == 1023){
42             answer(p[1], p[1024]);
43         }else if(N == 1535){
44             plot(512, 1024);
45             answer(p[1], p[1536]);
46         }else if(N == 1791){
47             plot(256, 1024);
48             answer(p[1], p[1792]);
49         }else if(N == 1919){
50             plot(128, 1024);
51             answer(p[1], p[1920]);
52         }else if(N == 1983){
53             plot(64, 1024);
54             answer(p[1], p[1984]);
55         }
56     }else{
57         for(int i = 0; (1 << i) < N; i++){
58             plot(c, c + (1 << i));
```

命令提示字元

```
solution.cpp:19:24: note: 'std::pair<int, int> X' previously declared here
19 |         pair<int, int> X = intersections(c1, c2);
    |         ^
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>grader_mul
OK
21
1535 1791 1919 1983
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>g++ grader_multiple.cpp solution.cpp -o grader_mul
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>grader_mul
WA
The distance between the two points is 6.0000000 instead of 1535.
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>g++ grader_multiple.cpp solution.cpp -o grader_mul
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>grader_mul
OK
20
767 895 959 991 1007 1015 1019 1021 1022 1279 1407 1471 1503 1519 1527 1531 1533 1534 1663 1727 1759 1775 1783 1787 1789 1
790 1855 1887 1903 1911 1915 1917 1918 1951 1967 1975 1979 1981 1982 1999
C:\Users\User\Documents\code\problemsetting\TW2221 - Typhoon Kompasu>
```

Compiler (17) Resources Compile Log Debug Find Results Close

Line	Col	File	Message
77	39	C:\Users\User\Documents\code\problemsetting\TW2221...	In member function 'point line::unit()':
77	39	C:\Users\User\Documents\code\problemsetting\TW2221...	[Error] expected ')' before ')' token
77	39	C:\Users\User\Documents\code\problemsetting\TW2221...	[Error] could not convert '((((line*)this)'
77	39	C:\Users\User\Documents\code\problemsetting\TW2221...	[Error] expected ';' before ')' token
77	39	C:\Users\User\Documents\code\problemsetting\TW2221...	[Error] expected primary-expression before ')' token
111	95	C:\Users\User\Documents\code\problemsetting\TW2221...	In function 'point intersection(line, line, int)':
111	95	C:\Users\User\Documents\code\problemsetting\TW2221...	[Error] 'WA' was not declared in this scope

Line: 54 Col: 32 Sel: 0 Lines: 66 Length: 1371 Insert

CH 速

26°C 晴時多雲

23:03 25/5/2022

啟用 Windows
移至 [設定] 以啟用 Windows。

Can it be further?

As you see there are a lot of equal cases

~~Hard code each then C can be reduced to 19~~

Too time consuming!

Time to change our approach

Is there any general method to tackle the case that there are many 'I's?

$$C = 15$$

If there are two '1' continuous, we move left by 1

We can ignore smaller bits recursively

$$C = 15$$

Score: 91.5

 命令提示字元

移至 [設定] 以啟用 Windows。

Line: 21 Col: 28 Sel: 0 Lines: 38 Length: 792 Insert

Final Countdown (C = 12)

Move left move right repeatedly

Theoretical maximum!

C = 12

Expected Score: 100

Circle-circle intersection is never used!

Indeed, this is just a camouflage since it is unnatural to 'ban' circle-circle intersection 😊

DVD Logo

TW2222

Problem Statement

A DVD Logo bouncing, starting at (x, y) and initially move (v_x, v_y) per second.

If it will eventually hit the corner, output **YES** on the first line and the time needed on the second line.

Otherwise, output **NO**.

Subtask I

$$v_x = 0$$
$$1 \leq N, M \leq 1000$$

Trivial case work.

- $y = 0$ or M $\rightarrow$ YES
- Otherwise $\rightarrow$ NO

Score: 7

Subtask 2

$$v_x = v_y = 1$$
$$1 \leq N, M \leq 1000$$

Simulate.

Remember to check first whether it is initially at corner.

Score: 30

Subtask 3

N and x is divisible by v_x

M and y is divisible by v_y

If v_x is negative, can set $v_x := -v_x$ and $x := n - x$ without changing the answer.

Similar for v_y .

Speed up simulation.

Score: 45

Partial Score - Yes/No only

Extend the TV

We have $N \mid x + tv_x$, $M \mid y + tv_y$, t is a rational number

So $aN = x + tv_x$, $bM = y + tv_y$

$t = (aN - x) / v_x$

$bM = y + (aN - x) / v_x * v_y$

$bM * v_x - aN * v_y = y * v_x - x * v_y$

So YES iff $(bv_x, av_y) \mid y * v_x - x * v_y$

Score: 35.6

Subtask 4-5

$$bM * v_x - aN * v_y = y * v_x - x * v_y$$

Seems complicated!

But it is just in form of

$$bX - aY = Z, X = Mv_x, Y = Nv_y, Z = yv_x - xv_y$$

Since X and Y are nonnegative integers, we may increase X and Y depending on $bX - aY > Z$ or $bX - aY < Z$

If $bX - aY > Z$, increase Y by 1

If $bX - aY < Z$, increase X by 1

If $bX - aY = Z$, we have find the solution: output $t = (aN - x) / v_x$

Repeatedly until X and Y are too large (e.g. $> 1e6$), and output 'NO'

Score: 83 (wrong initial: 28.8)

Subtask 6

Solve $bM * v_x - aN * v_y = y * v_x - x * v_y$ by extended euclidean algorithm

Note that $a \geq x / N$ and $b \geq y / M$

Solving it and we are done

Remember to compute division in floating point in the last step

Score: 100

Colourful Highlighter

TW2223

Problem Statement

There are N highlighters

You may do the two type of operations:

moveLeft(i): Move the highlighter currently indexed i to left end.

moveRight(i): Move the highlighter currently indexed i to right end.

Find the minimum number of operations needed, and output any one of the possible sequence of operations.

Subtask I

Hard code each of the 6 cases for $N = 3$

1 2 3

1 3 2

2 1 3

2 3 1

3 1 2

3 2 1

Score: 5

Subtask 2

Brute force simulation

Observation 1. You never move a highlighter more than once.

Therefore, you may just choose an order and 'L' or 'R' and simulate them.

Time Complexity: $O(N * N! * 2^N)$

Score: 17

Subtask 3

Define L and R the number of times `moveLeft()` and `moveRight()` is used respectively

Observation 2: `moveLeft()` and `moveRight()` does not affect each other

Observation 3a: If L and R is fixed, if c_i are all unique, then there exist only one possible choice to move the highlighters, from large to small **in value** for L and from small to large **in value** for R.

Just go through all pairs of L, R, and find the order

Time complexity: $O(N^3)$

Score: 28 (Cumulative: 40)

Subtask 4

~~Observation 3a: If L and R is fixed, if c_i are all unique, then there exist only one possible choice to move the highlighters, from large to small for L and from small to large for R.~~

No longer holds as there are equal values in c_i

Observation 3b: If L and R is fixed, there exist only one **best** choice to move the highlighters, from large to small **in value and index** for L and from small to large **in value and index** for R.

L increase $\rightarrow$ R increase, so just have to check $O(N)$ pairs of (L, R)

Remember subtask 1 cases

Time complexity: $O(N^2)$

Score: 53

Subtask 5

To maintain whether an array is sorted, we may notice that if it is sorted, then there exist a boundary for $C = 1$ and 2

We maintain the left and right bound of $C = 1$ and 2 respectively, and update them accordingly

$O(1)$ update when L or R change

After finding L and R , just choose the last L '1' and first R '2' and move them to the left and right respectively.

Time complexity: $O(N)$ or $O(N \lg N)$

Score: 30 (Cumulative: 83)

Subtask 6

Combine idea of subtask 4 and 5

Use a set to maintain the 'sorted' sequence in the middle

Increase R until can't, then increase L

Time Complexity: $O(N \lg N)$

Score: 100

Blossom Tree

TW2224

Problem Statement

東風夜放花千樹。更吹落、星如雨。
寶馬雕車香滿路。鳳簫聲動，玉壺光轉，一夜魚龍舞。
蛾兒雪柳黃金縷。笑語盈盈暗香去。
衆裏尋他千百度。驀然回首，那人卻在，燈火闌珊處。

《青玉案·元夕》

There are N trees in $[-P, P]$, each at position x_i with lightness of l_i

The contribution of light by tree i at position x is $\max(0, l_i - \text{abs}(x - x_i))$

Find the dimmest position and the respective lightness

Subtask I

For $N = 1$,

Check $x = -P$ and P . Find the dimmer one.

For $N = 2$,

Check $x = -P, P$ and start/end point of each light segment in the middle.

Always remember to find the one with the largest index

Expected score: 8

Subtask 2 (Solution 1)

Loop over N and P

For each position $i = -P, -P+1, \dots, P$:

$light = 0$

 for each tree j :

$light += l_j - \text{abs}(i - x_j);$

 update ans

Time complexity: $O(NP)$

Expected Score: 12

Subtask 2 (Solution 2)

Loop over N and L instead

For each tree i :

```
for (j = max(P, x_i - l_i); j <= min(P, x_i + l_i); j++):  
    light[j] += l_j - abs(i - x_j)
```

check ans

Score: 12 (Cumulative: 20)

Subtask 3

Special points

Note that only the points $-P, P, x[i] - l[i], x[i], x[i] + l[i]$ are 'special'

For any other points, we can move left or move right to further minimize the lightness

[add graph]

There are $O(3N)$ special points

Time complexity: $O(3N) * O(N) = O(N^2)$

Score: 30 (Cumulative: 50)

Subtask 4

Partial-partial sum

Or just sort the indexes and maintain it on the go

[Insert diagram here]

Score: 45 (Cumulative: 83)

Subtask 5

Combine idea of subtask 3 and 4

We may maintain the double partial difference and compute back the lightness of all the special points

Score: 100